

P-15944

SUPERINTENDENCIA DE INDUSTRIA Y COMERCIO

DIRECCIÓN DE NUEVAS CREACIONES

E. S. D.

Re: Expediente No. **NC2018/0009316**

Solicitud de patente a nombre de **CRYPTZONE NORTH AMERICA INC.**

Yo, CAROLINA DAZA MONTALVO identificada como aparece al pie de mi firma, domiciliada en la Carrera 11 No. 86-53, obrando como apoderada de la sociedad CRYPTZONE NORTH AMERICA INC., domiciliada en Massachusetts, Estados Unidos de América, solicitante de la patente de la referencia, en respuesta al oficio No. **7273** notificado electrónicamente el 30 de julio de 2019, relacionado con el concepto técnico de fondo, me permito manifestar:

1. Modificación a las Reivindicaciones

Con fundamento en la facultad que confiere el artículo 34 de la Decisión 486 al solicitante para modificar la solicitud de patente en cualquier momento del trámite, presento junto con este memorial nuevo capítulo reivindicatorio, el cual consta de veinte (20) reivindicaciones y reemplaza completamente el capítulo presentado anteriormente. En el nuevo capítulo reivindicatorio se realizaron las siguientes modificaciones:

- 1.1** Se modificó la antigua reivindicación 1, para recitar ahora que: *"en donde la red privada comprende uno o más dispositivos de red; recibir, por el primer sistema informático, desde el dispositivo del cliente, una lista de acceso del cliente que indica las de los dispositivos de red en la red privada que pueden comunicarse con el dispositivo del cliente", "iniciar, para el primer túnel de red" y "en donde cada una de las reglas del cortafuegos se derivan de la lista de acceso del cliente", y así dar mayor claridad al método reclamado.*
- 1.2** De forma similar, se modificó la antigua reivindicación 11, para recitar ahora que: *"en donde la red privada comprende uno o más dispositivos de red; recibir, del dispositivo del cliente, una lista de acceso del cliente que indica las de los*

dispositivos de red en la red privada que pueden comunicarse con el dispositivo del cliente”, “iniciar, para el primer túnel de red” y “en donde cada una de las reglas del cortafuegos se derivan de la lista de acceso del cliente”, y así dar mayor claridad al sistema reclamado.

- 1.3** De igual manera, se modificó la antigua reivindicación 16, para ahora recitar: *“recibir, por el primer sistema informático, del dispositivo del cliente, una lista de acceso del cliente”, “iniciar, para el primer túnel de red” y “en donde cada una de las reglas del cortafuegos se derivan de la lista de acceso del cliente”, y así dar mayor claridad al medio de almacenamiento reclamado.*

En cuanto a estas modificaciones, me permito decir que las mismas no amplían de ninguna manera la materia objeto de la presente solicitud y se realizaron con el ánimo de superar las objeciones de patentabilidad planteadas por el examinador en su concepto técnico.

2. Título

Es deseo del solicitante manifestar a este Despacho que se realiza el cambio del título de la invención a *“MÉTODO, SISTEMA Y MEDIO DE ALMACENAMIENTO LEGIBLE POR COMPUTADOR PARA PROTECCIÓN DE DISPOSITIVOS DE RED MEDIANTE UN CORTAFUEGOS”*, como lo ha sugerido el examinador.

3. Pago de tasas

En relación con el número de reivindicaciones, atentamente nos permitimos decir que no se adjunta tasa por reivindicaciones adicionales, toda vez que el nuevo capítulo reivindicatorio contiene el mismo número de cláusulas que el anteriormente presentado, y en esencia es el mismo objeto, pero definido con mayor claridad.

Por otra parte, considerando que el presente memorial se presenta junto con modificaciones a las reivindicaciones, se adjunta recibo por concepto de examen de patentabilidad, según lo establecido en la Resolución 3719 del 2016.

3. Novedad

El Examinador establece en su concepto técnico la falta de novedad de las antiguas reivindicaciones 1 a 16 y 11 a 20 a la luz de las enseñanzas técnicas descritas en los documentos US20150293756 (D1) y US20040225895 (D2).

A este respecto, me permito decir que el artículo 16 de la Decisión 486 de la Comisión de la Comunidad Andina establece que la invención se considera nueva cuando no está comprendida en el estado de la técnica. Si de acuerdo con el espíritu de la norma legal indicada, la invención no se considera como nueva si está "comprendida" en el estado de la técnica, dicha norma se está refiriendo, naturalmente, a la totalidad del invento y no a alguno o a algunos de los componentes.

Aplicando lo anterior al caso concreto, para que las razones que aduce ese Despacho para desvirtuar la novedad fuesen ciertas, sería forzoso que se demostrara que la invención de mi mandante esté comprendida en el estado de la técnica, específicamente en los documentos D1 y D2, de forma TOTAL, lo cual no sucede con la materia reclamada en el nuevo capítulo reivindicatorio.

Considerando las anteriores definiciones y teniendo en cuenta las modificaciones realizadas al capítulo reivindicatorio, es deseo de la solicitante manifestar su desacuerdo con dichas objeciones y presentar argumentos que demuestran la novedad de la invención que se reclama en las nuevas reivindicaciones 1 a 20.

En primer lugar, se llama la atención del examinador sobre la materia reclamada en el Capítulo Reivindicatorio, en el cual, las nuevas reivindicaciones independientes 1, 11 y 16 reclaman de forma similar **"... en donde la red privada comprende uno o más dispositivos de red; recibir, por el primer sistema informático, desde el dispositivo del cliente, una lista de acceso del cliente que indica las de los dispositivos de red en la red privada que pueden comunicarse con el dispositivo del cliente; y iniciar, para el primer túnel de red, un servicio de cortafuegos separado con un conjunto separado de reglas de cortafuegos en el primer sistema informático para bloquear selectivamente y permitir el tráfico de red entre el dispositivo del cliente y el uno o más dispositivos de red en la red privada, en donde cada una de las reglas del cortafuegos se derivan de la lista de acceso del cliente"**. Tal como puede establecer una persona del oficio normalmente versada en la materia, ninguna de estas características técnicas está presente en los documentos D1 y D2, citados por el examinador.

El solicitante considera que el objeto de las reivindicaciones independientes 1, 11 y 16, es por tanto novedoso respecto al estado de la técnica por las siguientes razones:

Una objeción de novedad requiere que una sola referencia describa de manera idéntica todas y cada una de las características de la reivindicación independiente 1 del solicitante. La reivindicación 1 se ha modificado anteriormente para ahora recitar la recepción de una lista de acceso del cliente desde un dispositivo cliente y derivar las reglas del cortafuegos de la lista de acceso del cliente.

Al objetar la reivindicación dependiente 6, que se relaciona con la recepción de una lista de acceso del cliente y la determinación de las reglas del cortafuegos de la lista de acceso del cliente, el examinador cita a D1 en los párrafos [0017], [0055] y [0081]. Al menos por las siguientes razones, D1 no enseña o sugiere las características anteriores de la reivindicación 1.

En particular, el párrafo [0017] de D1 describe el establecimiento de un túnel entre un nodo y un nodo maestro. Aunque se menciona una conexión de red privada virtual, **no se enseña recibir una lista de acceso de cliente desde un dispositivo cliente**, ni a derivar reglas del cortafuegos de esta lista de acceso de cliente.

El párrafo [0055] de D1 describe **el monitoreo utilizando un servidor centralizado** que envía solicitudes de determinación de estado a máquinas virtuales en dispositivos para obtener información de estado. Sin embargo, estas solicitudes para obtener información del estado **no enseñan ni sugieren recibir una lista de acceso del cliente desde un dispositivo cliente y usar esa lista para derivar las reglas del cortafuegos**, como se reclama.

Finalmente, el párrafo [0081] de D1 simplemente describe una segunda máquina virtual a la que se le da la responsabilidad de proteger el tráfico de la red y que es controlada por un administrador del sistema central. Pero este control **no enseña ni sugiere recibir reglas de acceso del cliente desde un dispositivo cliente**, como se reclama y se discutió anteriormente.

Otros párrafos de D1 citados por el examinador respetuosamente no remedian las deficiencias anteriores. Por ejemplo, el párrafo [0072] de D1 describe un maestro tunelizador que forma parte del sistema de administración **central** e instruye a una puerta de enlace para establecer una conexión. Sin embargo, este maestro de túnel **no**

enseña ni sugiere que se establezca un primer túnel de red en respuesta a la solicitud de un dispositivo cliente, que este dispositivo cliente proporcione una lista de acceso de cliente y que las reglas del cortafuegos se deriven de esta lista de acceso de cliente para la primera red túnel, como se reclama.

El párrafo [0073] de D1 también describe el sistema de gestión **central**. En particular, D1 describe que este sistema de administración central brinda detalles de rutas y qué máquinas virtuales están hablando entre sí. Sin embargo, este sistema de administración central **no enseña ni sugiere** recibir una lista de acceso de clientes desde un dispositivo cliente, como se reclama y se discutió anteriormente.

Por consiguiente, al menos por las razones discutidas anteriormente, la reivindicación independiente 1 es patentable. En cuanto a las reivindicaciones independientes 11 y 16 del solicitante son patentables al menos por razones similares como se discutió anteriormente para la reivindicación independiente 1.

Adicional a lo anterior, es deseo del solicitante indicar que:

Si bien entendemos y aceptamos que cada país es autónomo en la concesión o negación de los privilegios de patente, queremos resaltar el hecho que el análisis de los requisitos para el otorgamiento de esos derechos es prácticamente universal. La presente invención fue concedida por la Oficina de Patentes de los Estados Unidos bajo el número US 9,628,444, indica que las reivindicaciones con el mismo alcance son novedosas e inventivas sobre la técnica anterior. Se presenta como anexo al presente memorial copia de la Patente No. US 9,628,444.

Por lo tanto, claramente existe por lo menos una diferencia entre la presente invención definida en las reivindicaciones independientes y los documentos D1 y D2, con lo cual se ajusta a la definición de Novedad, como se indicó anteriormente, razón por la cual dichas anterioridades no desvirtúan dicho requisito de forma completa.

Con base en el análisis anterior, D1 y D2 no enseñan o sugieren la presente invención. En vista de lo anterior, **las reivindicaciones presentadas son novedosas e inventivas sobre el estado del arte**. Por lo tanto, se le solicita respetuosamente al señor examinador retire la objeción de novedad.

De esta manera, queda demostrado que el documento D1 citado por su Despacho **no puede ser considerado como anterioridad válida** capaz de afectar la novedad de la solicitud, toda vez que este documento falla en enseñar todas las características esenciales de la invención, y la manera como se solucionó el problema técnico en cuestión.

En este sentido, y siguiendo el Manual Andino de Patentes, la presente invención tal y como se encuentra definida en las reivindicaciones 1 a 20, tiene diferencias con el estado del arte e implicó un nivel inventivo durante su desarrollo, por lo cual debe ser considerada como novedosa e inventiva según las disposiciones de los artículos 16 y 18 de la Decisión 486.

En consecuencia, solicitamos a ese despacho pronunciarse en relación con todos los argumentos presentados en este memorial dando estricta aplicación al principio del debido proceso consagrado en el artículo 29 de la Constitución Nacional.

5. Anexos

- Capítulo reivindicatorio que incluye veinte (20) reivindicaciones.
- Copia de la Patente No. US 9,628,444.
- Recibo de pago por concepto de examen de patentabilidad.

Señora Directora,



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(12) **United States Patent**
Glazemakers et al.

(10) **Patent No.:** **US 9,628,444 B1**
(45) **Date of Patent:** **Apr. 18, 2017**

(54) **PROTECTING NETWORK DEVICES BY A FIREWALL**

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(*) Notice: Subject to any disclaimer, the term of this
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(21) Appl. No.: **15/053,422**

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Primary Examiner — Christopher Revak

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Related U.S. Application Data

(60) Provisional application No. 62/292,702, filed on Feb.
8, 2016.

(51) **Int. Cl.**
H04L 29/06 (2006.01)

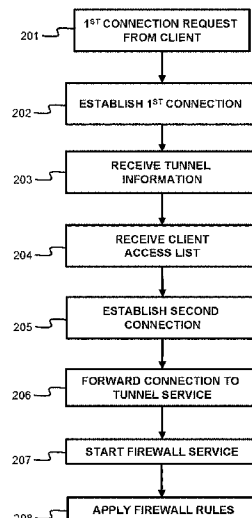
(52) **U.S. Cl.**
CPC **H04L 63/029** (2013.01); **H04L 63/0254**
(2013.01); **H04L 63/0272** (2013.01); **H04L**
63/0281 (2013.01)

(58) **Field of Classification Search**
None
See application file for complete search history.

(57) **ABSTRACT**

Systems and methods provide for scaling and management of a gateway. In one embodiment, a method includes: in response to a request from a client device, establishing, by a computer system implementing a gateway to a private network, a network tunnel between the client device and the gateway; and after establishing the network tunnel, starting a separate firewall service with a separate set of firewall rules on the computer system for selectively blocking and allowing network traffic between the client device and one or more network devices in the private network.

20 Claims, 6 Drawing Sheets



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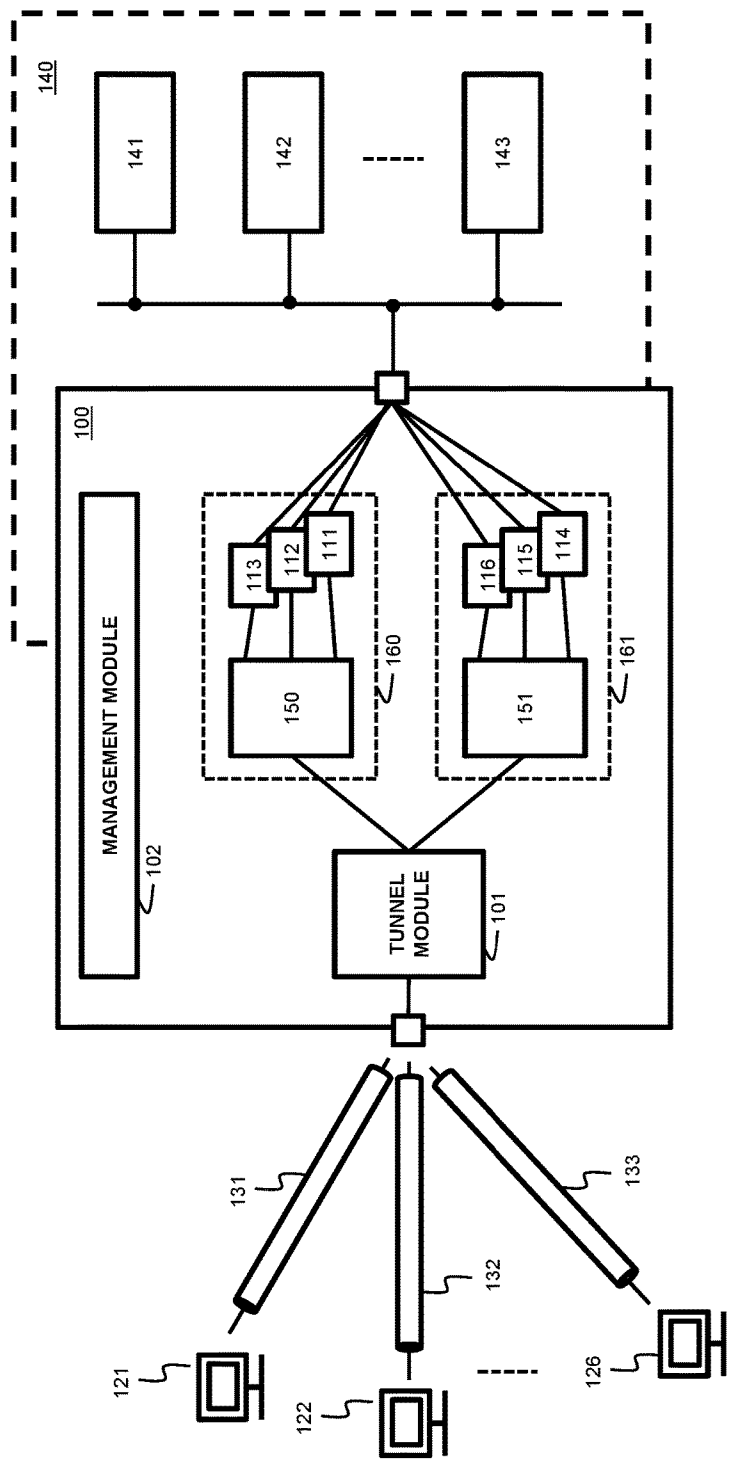


Fig. 1

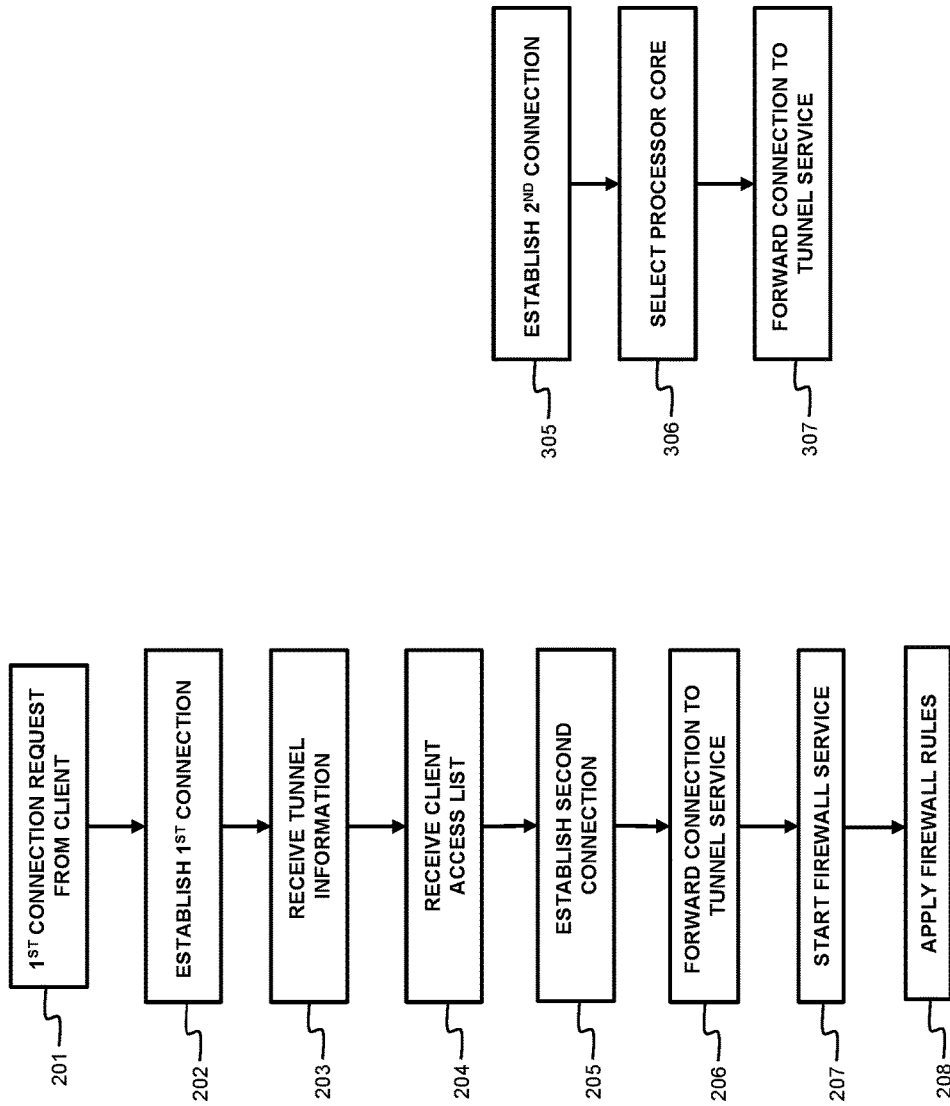


Fig. 3

Fig. 2

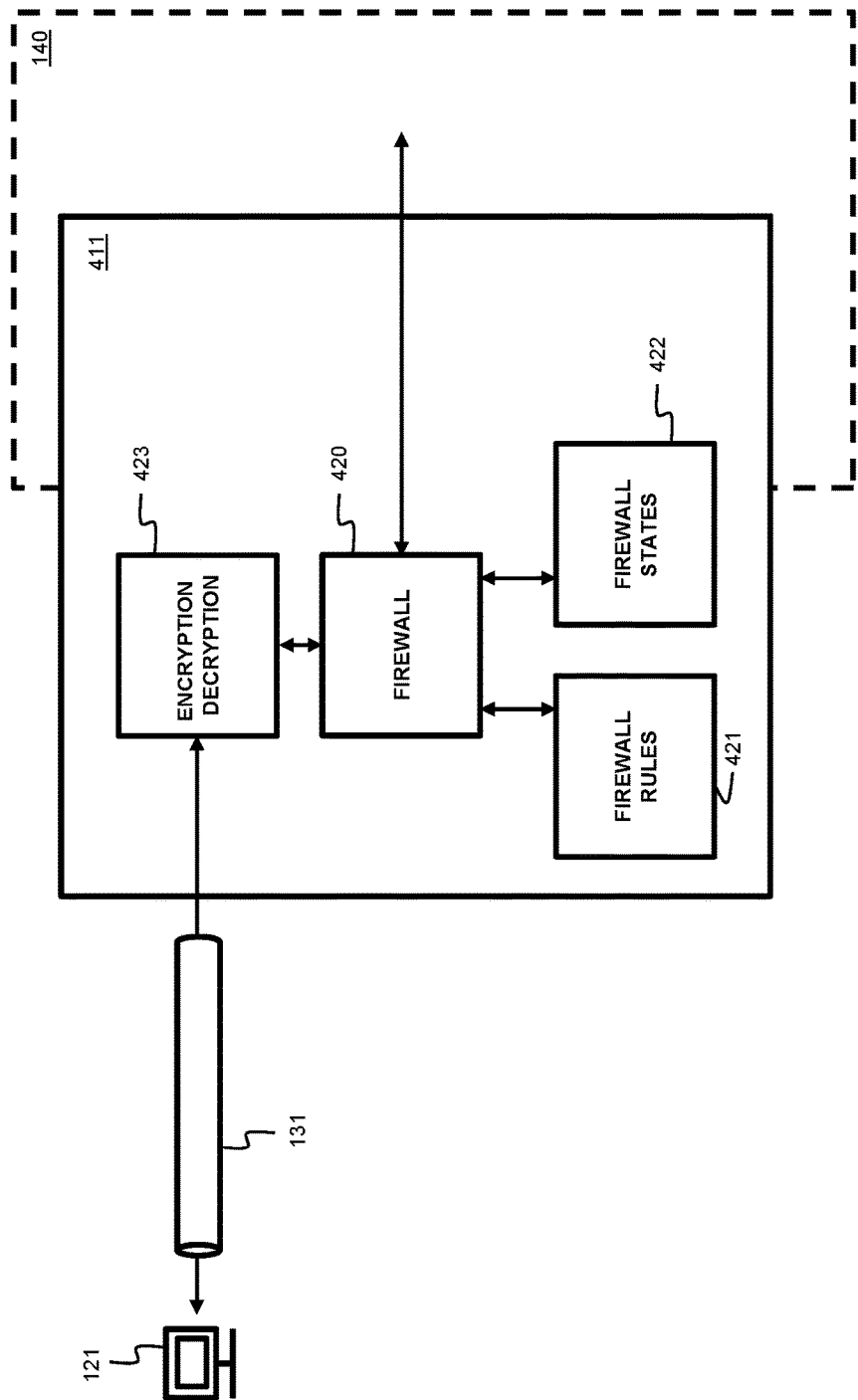


Fig. 4

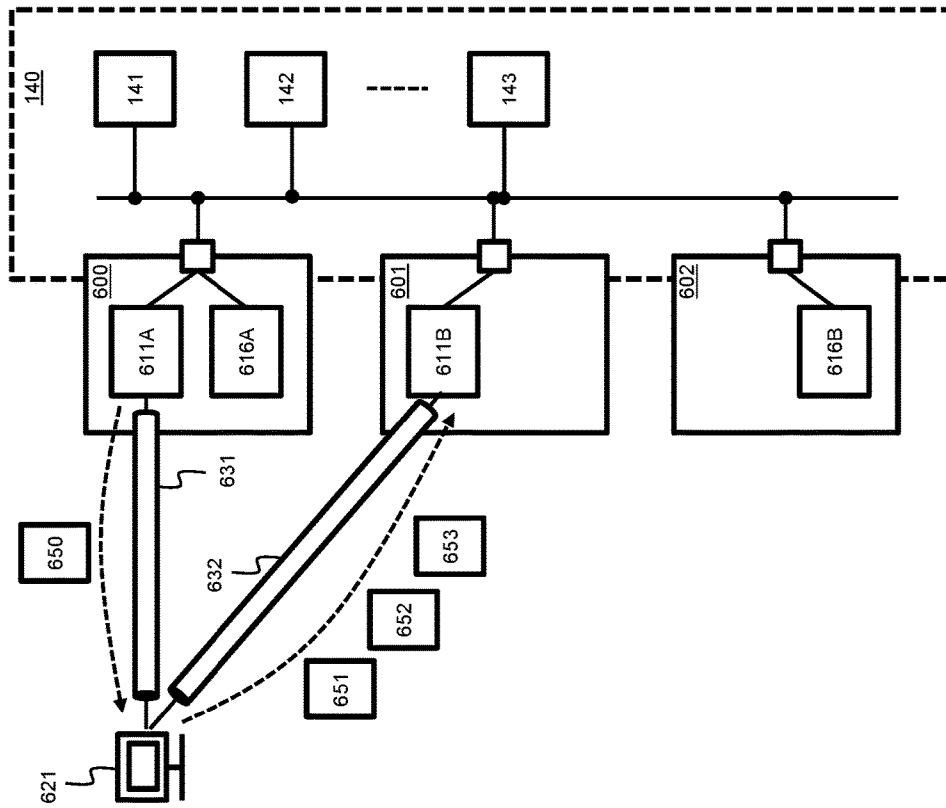


Fig. 6

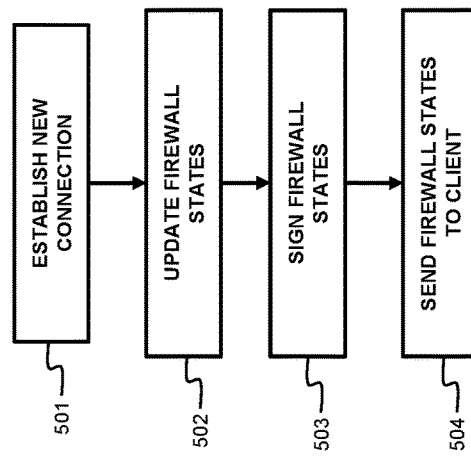


Fig. 5

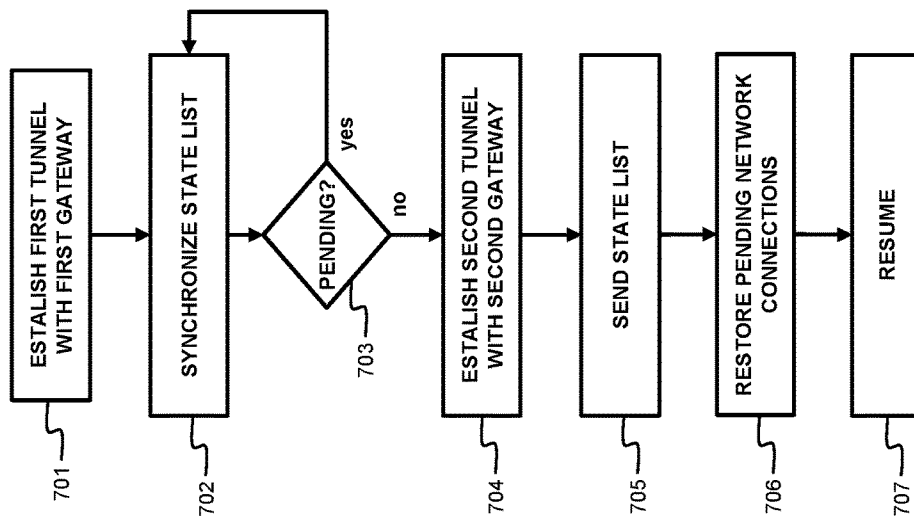


Fig. 7

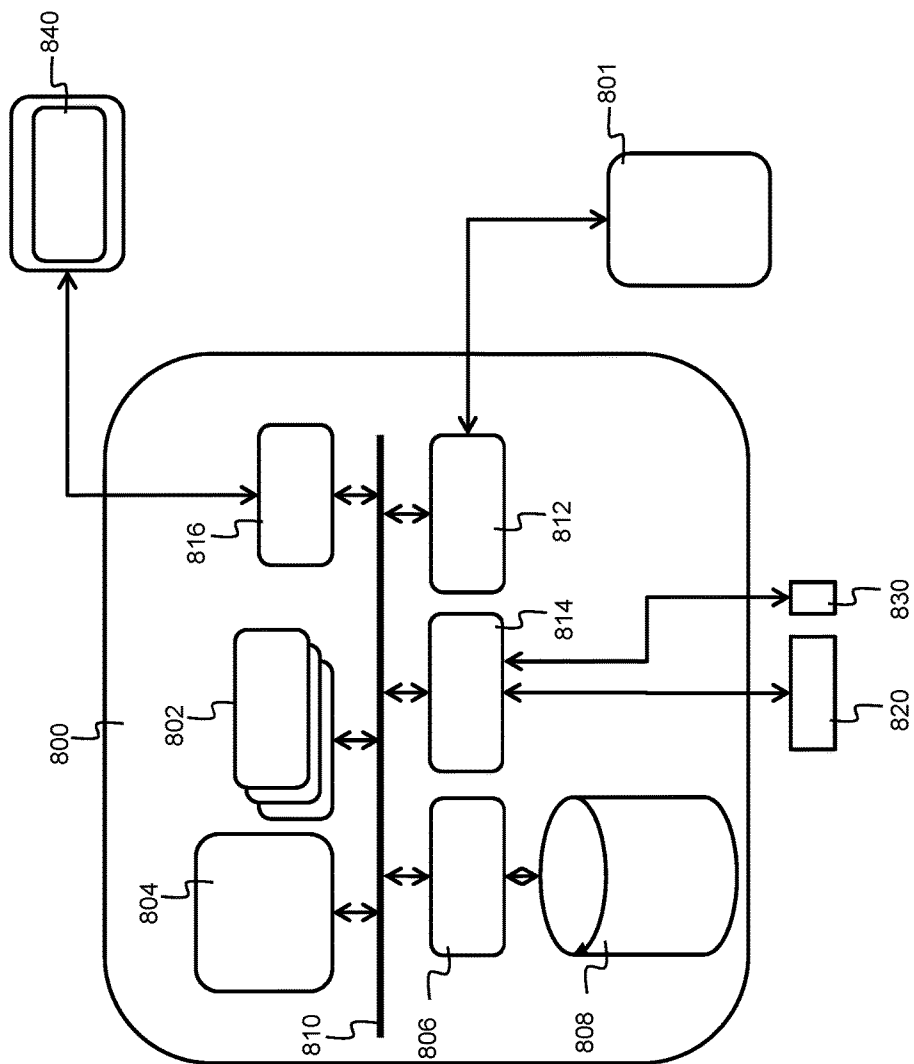


Fig. 8

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PROTECTING NETWORK DEVICES BY A FIREWALL

RELATED APPLICATIONS

This application claims the benefit of and priority to U.S. Provisional Application Ser. No. 62/292,702, filed Feb. 8, 2016, entitled "SYSTEMS AND METHODS FOR PROTECTING NETWORK DEVICES BY A FIREWALL," by Glazemakers et al., the entire contents of which application is incorporated by reference as if fully set forth herein.

FIELD

At least some embodiments generally relate to the field of network protection and more particularly, but not limited to, the protection of private networks by a gateway including a tunnel server such as a VPN server and a firewall.

BACKGROUND

In order to protect private networks from unwanted network access, a firewall may be implemented in a gateway in order to selectively filter communication from and to the private network. By applying firewall rules, the firewall then lets network packets pass, or blocks them in one or both directions. The rules are typically defined by the source and/or destination addresses of the network packets, or the source and/or destination ports of the network packets.

A firewall may further perform stateful packet inspection thereby keeping track of the state of networking connections such as, for example, TCP or UDP networking connections. This way, a finer-grained control is possible as the rules become dependent on the network connection and thus dynamic. The firewall keeps track of all open network connections by maintaining a table, the state table or state list, with all the pending connections.

In order to have better control on the network traffic passing by a gateway, the firewall may further perform deep packet inspection by controlling the payload data of the networking packets. This way, different types of control can be implemented such as, for example, user access or authorization, protocol compliance, spam filtering and virus detection.

In order to protect a private network further, the gateway's firewall may be combined with network tunnelling. Access to the private network may then be established by a VPN (Virtual Private Network) where a secured networking tunnel is set up between a client device and the gateway. The setup of such a tunnel is only granted upon successful authentication with the gateway, which then functions as a VPN server. By the combination of a firewall and VPN server in the gateway, access to devices in the private network can be authorized on a client or a user level by the VPN server and on a network level by the firewall.

SUMMARY

Several problems may arise with combined gateway systems as described above. First, the amount of firewall rules will increase with each additional user, making the management of all the rules difficult for large systems. Because of this, each incoming packet will have to be processed against all of these rules.

A second problem relates to the failover of the gateway. In one failover scheme, referred to as passive failover, all the information in the gateway (i.e., firewall rules, state list and

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active networking tunnels) is constantly synchronized with a second redundant gateway. If there is a failure of the first gateway, the second gateway automatically takes over from the first gateway. Because of the synchronization, no active network connections are lost during failover.

In another failover scheme, referred to as active failover, several gateways are also synchronized with each other, but all are active at the same time such that the load can be balanced between the gateways. Also, in this case, because of the synchronization, one gateway can take over all connections from another gateway upon failure without losing the network connections.

A problem with the two schemes is scalability because the system cannot be scaled by just adding new gateways. First, the bandwidth between the gateways needed for synchronization will increase exponentially with the amount or number of gateways. Second, because the state list is kept synchronized, the size of the state list and firewall rules for each gateway will also increase linearly when scaling. Because of this, the needed bandwidth, processing power and memory space of each gateway will increase when adding a new gateway.

Various embodiments are described herein that alleviate the above problems and provide a gateway that can be scaled and managed in an easy and straightforward manner. Some embodiments are summarized in this section.

Embodiments of the present disclosure help protect network devices from unauthorized access. Among other things, embodiments of the disclosure allow full access to application servers and other network devices that a client is authorized to access, while preventing all access (or even knowledge) of network devices the client is not authorized to access.

A computer-implemented method according to one embodiment of the present disclosure includes: upon a request from a client device, establishing, by a computer system implementing a gateway to a private network, a network tunnel between the client device and the gateway; and, upon establishing the network tunnel, starting a separate firewall service with a separate set of firewall rules on the computer system for selectively blocking and allowing network traffic between the client device and one or more network devices in the private network.

The present disclosure includes various methods, apparatuses (including computer systems) that perform such methods, and computer readable media containing instructions that, when executed by computing systems, cause the computing systems to perform such methods.

Other features will be apparent from the accompanying drawings and from the detailed description which follows.

BRIEF DESCRIPTION OF THE DRAWINGS

The embodiments are illustrated by way of example and not limitation in the figures of the accompanying drawings in which like references indicate similar elements.

FIG. 1 is an exemplary system for the protection of network devices from unwanted network access according to various aspects of the present disclosure.

FIG. 2 is an exemplary process for establishing a networking tunnel between a client device and a private network that may be executed by components of the present disclosure.

FIG. 3 is an exemplary process for distributing firewall services among different processor cores of a computing system implementing a gateway according to one embodiment.

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FIG. 4 is an exemplary system for the establishment of a networking tunnel between a client device and a private network separated by a gateway implementing a firewall service according to one embodiment.

FIG. 5 is an exemplary process for synchronizing a firewall implemented on a gateway that may be executed by components of the present disclosure.

FIG. 6 is an exemplary system including a plurality of gateways for the protection of network devices from unwanted network access according to various aspects of the present disclosure.

FIG. 7 is an exemplary process for seamlessly transferring a firewall service from one gateway to another gateway executed by components of the present disclosure.

FIG. 8 illustrates an exemplary computing system according to various aspects of the present disclosure.

DETAILED DESCRIPTION

Subject matter will now be described more fully herein after with reference to the accompanying drawings, which form a part hereof, and which show, by way of illustration, specific example embodiments. Subject matter may, however, be embodied in a variety of different forms and, therefore, covered or claimed subject matter is intended to be construed as not being limited to any example embodiments set forth herein; example embodiments are provided merely to be illustrative. Likewise, a reasonably broad scope for claimed or covered subject matter is intended. Among other things, for example, subject matter may be embodied as methods, devices, components, or systems. Accordingly, embodiments may, for example, take the form of hardware, software, firmware or any combination thereof (other than software per se). The following detailed description is, therefore, not intended to be taken in a limiting sense.

In the accompanying drawings, some features may be exaggerated to show details of particular components (and any size, material and similar details shown in the figures are intended to be illustrative and not restrictive). Therefore, specific structural and functional details disclosed herein are not to be interpreted as limiting, but merely as a representative basis for teaching one skilled in the art to variously employ the disclosed embodiments.

Reference in this specification to “one embodiment” or “an embodiment” means that a particular feature, structure, or characteristic described in connection with the embodiment is included in at least one embodiment of the disclosure. The appearances of the phrase “in one embodiment” in various places in the specification are not necessarily all referring to the same embodiment, nor are separate or alternative embodiments mutually exclusive of other embodiments. Moreover, various features are described which may be exhibited by some embodiments and not by others. Similarly, various requirements are described which may be requirements for some embodiments but not other embodiments.

Any combination and/or subset of the elements of the methods depicted herein may be combined with each other, selectively performed or not performed based on various conditions, repeated any desired number of times, and practiced in any suitable order and in conjunction with any suitable system, device, and/or process. The methods described and depicted herein can be implemented in any suitable manner, such as through software operating on one or more computer systems. The software may comprise computer-readable instructions stored in a tangible computer-readable medium (such as the memory of a computer

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system) and can be executed by one or more processors to perform the methods of various embodiments.

FIG. 1 illustrates an exemplary system for the protection of network devices from unwanted network access according to various aspects of the present disclosure. In this example, three network devices (application servers **141**, **142** and **143**) are part of a private network **140**. Access to the servers **141-143** is obtained from within the private network **140** via a private network address. In this context, the term “private” refers to the fact that the application servers **141-143** are not globally routable. In other words, the application servers **141-143** cannot be addressed by their private network address from outside the private network **140**.

The private network **140** and other components in FIG. 1 may utilize any number and type of communication protocols, also referred to as the Internet Protocol (“IP”), or as the Transmission Control Protocol/Internet Protocol (“TCP/IP”). For example, the private network **140** may have address ranges as set by RFC 1918 for Internet Protocol Version 4 or IPv4 and RFC 4193 for Internet Protocol Version 6 or IPv6.

Networking devices **141-143** may correspond to application servers that provide services over the network **140** to other computing devices. Any number and type of application servers and associated services may be used in conjunction with embodiments of the present disclosure, such as mail servers, file servers, Customer Relationship Management or CRM services, Enterprise Resource Planning or ERP services, and/or document management services.

A data connection may then be established with any of application servers **141-143** by opening a communication socket with the respective application server on a port (or port range) associated with the service. Application servers **141-143** may correspond to physical devices with a physical networking interface associated with a private network address. Alternatively, application servers **141-143** may also correspond to virtual server instances running on one or more physical servers. Virtual server instances may each have a virtual network interface with an associated private network address. Virtual server instances may include, as well as operate in conjunction with, one or more user space instances (also known as software containers, virtual engines, virtual private servers, and/or jails). Such user space instances may be implemented in any suitable manner, including, for example, via the DOCKER software tool.

In the example shown in FIG. 1, the private network **140** is separated from an external network by a gateway **100**, thereby allowing networking traffic between the external network and the private network **140** in a controlled way. The system of FIG. 1 can identify clients **121-126** as “trusted clients” having access rights to one or more of the application servers **141-143** within the private network **140** in order to use the services running thereon. The clients **121-126** may be, or include, physical hardware and/or virtual components. For example, a client **121-126** may include a virtual operating system running on a physical device, such as a mobile device. The system can also grant network access to a selection of the application servers **141-143** which the clients **121-126** are allowed to access, and deny network access to any application server the clients **121-126** are not allowed to access.

In order to control access by the clients **121-126** to the application servers **141-143**, networking tunnels **131-133** are established between the clients **121-126** and the gateway **100**. This way, the private network **140** is extended to the clients **121-126**. In some embodiments, a virtual private

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network (or “VPN”) is established through tunnels **131-133**. In this manner, a client **121-126**, although not physically in the private network **140**, is provided a private network address in the range of the private network **140**, and can thus potentially access all application servers **141-143** by their respective private network address (provided access is allowed, as is discussed in more detail below).

All initial network connection requests from clients **121-126** are handled by tunnel module **101** implemented in gateway **100**. In one embodiment, tunnel module **101** looks at the initial connection (e.g., typically a first TLS packet that is arriving) and then determines whether to handover the connection to management module **102** (discussed below) in the case of metadata, or to one of tunnel services **150, 151** (discussed below) in the case of the actual tunnel traffic (i.e., data) from the clients. Once the connection is handed over (e.g., this can be done by passing over a file descriptor to management module **102** or one of the tunnel service modules), tunnel module **101** is not involved anymore. Then, the connection goes directly from the client to management module **102**, or to the selected tunnel service **150, 151**.

The initial network connection request may be a connection request for establishing a new tunnel by a new client device or a new network connection request within an existing tunnel. The data travelling over the connections in the tunnels **131-133** may further be protected by encryption, such as according to the Internet Protocol Security (or “IPsec protocol,”) Transport Layer Security (or “TLS”) and/or Datagram Transport Layer Security (or “DTLS”). In an example, the tunnel module **101** operates according to TLS or SSL and sets up the networking connections as TCP network connections. In order to do so, the clients send the request to an open port or port range of the gateway **100**, preferably the standard port 443 for TLS/SSL encrypted TCP connections.

Gateway **100** further implements firewall services **111-116**. Each respective firewall service then implements a firewall for selectively blocking and allowing network traffic between the respective client device and the network devices **141-143** in the private network. Each firewall service also comprises a set of firewall rules defining the access rules for a respective client device. In other words, gateway **100** runs a separate firewall for each connected client **121-126**. An advantage of this is that the size of the firewall rules of a firewall service does not grow with the amount of connected clients. This way, an increase in connected clients does not result in a performance loss due to an increase in the amount of firewall rules that have to be managed by a single firewall service.

FIG. 2 shows an exemplary process that may be executed by components of the present disclosure including the gateway **100** according to FIG. 1. By this process, a client device **121-126** is connected to the private network **140** through the gateway **100**. As an example, the process will be described with reference to client device **121**. In a first step **201**, the tunnel module **101** receives a first connection request from client networking device **121** to establish a first network connection with the gateway **100**, for example a secured TCP network connection request received on port 443 of the gateway. Thereupon, the network connection is established under step **202**, for example, by a three-way handshake in the case of a TCP network connection. This first network connection is used to exchange control information between the client **121** and the gateway **100**, and more particularly with the management module **102** implemented in the gateway **100**. In order to know that the

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connection is for control purposes, the tunnel module may inspect the first data packet exchanged over each newly-established network connection. If the data packet is a control data packet, the tunnel module **101** identifies the network connection as a control connection and will redirect all further packets received over this connection to the management module **102**. The control data packet may, for example, be identified by inspecting a specific TLS extension field in the header of the TLS packet.

In the next step **203**, the management module **102** receives tunnel information from the client device **121** over the first network connection. This information is further referred to as the client tunnel list. This client tunnel list includes information in order to establish the networking tunnel **131** with the gateway such as, for example, authentication information for authenticating the client with the gateway. After successful authentication by the management module **102**, the process proceeds to step **204**.

In step **204**, the management module **102** receives a client access list from client **121**. This client access list comprises a listing of the networking devices or applications in the private network **140** that the client device is allowed to access. The listing may, for example, comprise a network address such as an IPv4 or IPv6 network address identifying the respective networking device in the private network **140**. Also, a port number or port number range may be specified in order to limit the access to the respective networking device to one or more specific applications. The listing may also identify the respective networking device by its host name or by a fully qualified domain name (FQDN). The listing may also identify the respective networking device indirectly by specifying a further service from where one or more respective networking devices may be received. Such a service may, for example, be the Amazon Web Services which enables retrieval of a list of network devices using the AWS REST API according to metadata assigned to virtual instances in the Amazon cloud.

The client **121** may further retrieve this client access list and/or tunnel list from an authentication service that manages the client access to private network **140**. In order to avoid a situation in which the client compromises the client access list and/or tunnel list, the lists may be made unalterable (i.e., protected such that an alteration of one or both of the lists may be detected by the management module **102**).

Various methods and systems for authenticating the client and providing a client and tunnel list are further disclosed in U.S. Pat. No. 9,148,408 B1, which is hereby incorporated by reference in its entirety.

Then, in step **205**, the tunnel module **101** establishes a second network connection with the client **121** upon request of the client **121**. As this is a new connection, the tunnel module **101** inspects the first data packet received over this second connection. This second connection is used for the actual networking tunnel **131** because the client **121** has already established the first connection for the exchange of control information with the gateway **100**. The inspected data packet is therefore identified as a first packet for the not yet established networking tunnel **131**.

More specifically, in one embodiment, two connections are required for each tunnel (**131,132,133**) (e.g., two TCP connections in the case of a TLS tunnel). One connection is for uploading the tokens (metadata), and the other connection is for the actual tunnel traffic. The tunnel **131** is only established after step **207** (see below), and is only allowing traffic after step **208** (see below), because the firewall is blocked by default.

Thereupon, in step 206, the tunnel module 101 passes the second network connection (the connection that is responsible for tunnelling the traffic coming from the client 121) to tunnel service 150. Tunnel services 150 and 151 handle the tunnel traffic (131,132,133). Tunnel services 150, 151 are responsible to decrypt/encrypt the traffic coming and going to the respective client (121, 122, . . . 126) implemented on the gateway 100. The tunnel service then verifies with the management module 102 whether the networking tunnel 131 can be established. In the following step 207, the tunnel service 150 establishes the networking tunnel 131 by starting the firewall service 111. From this moment on, all data exchanged over the second network connection and thus over the networking tunnel 131 passes through firewall service 111. In other words, payload data from the networking tunnel 131 is forwarded to the started firewall service 111. This firewall service 111 implements a firewall that blocks by default all traffic between the client 121 and private network 140. In the next step 208, firewall service 111 applies the appropriate firewall rules to the firewall by retrieving the firewall rules from the management module 102. The management module 102 on its turn derives these firewall rules from the client access list.

More specifically, in one embodiment, the access rules are basically descriptive firewall rules, and the management module 102 fills in some of these descriptions. Some examples are as follows:

- Allow TCP traffic to 1.1.1.1 port 80 (note that this is an access rule that is an exact match with a firewall rule).
- Allow TCP traffic to www.google.com port 80 (note that this is an access rule that requires translation to multiple FW rules such as:
 - Allow TCP traffic to 173.194.71.103 port 80
 - Allow TCP traffic to 173.194.71.105 port 80
 - Allow TCP traffic to 173.194.71.99 port 80
 - . . .)

Thus, one access rule can, for example, result in many firewall rules. In some cases, an access rule can be a call to an IAAS provider like AWS web services and allow TCP traffic to, for example the following: [all instances that have finance in their description].

By the process according to FIG. 2, a network tunnel 131 between the client device 121 and the gateway 100 is thus established together with a separate firewall with a separate set of firewall rules for selectively blocking and allowing network traffic between the client device 121 and the network devices 141-143 in the private network 140. This process is performed for every client 121-126 that connects to the private network 140 thereby obtaining the respective firewall services 111-116.

According to a further embodiment, the gateway 100 is implemented on a computer system comprising multiple processor cores. Management module 102, tunnel module 101, tunnel service 150 and firewall services 111, 112 and 113 may then be implemented on the computer system by computer-executable instructions executed on the processor cores. The gateway 100 then includes a separate tunnel service running on each processor core. This is illustrated by the exemplary embodiment of FIG. 1 where tunnel service 150 and firewall services 111, 112 and 113 are implemented on the first processor core 160 and tunnel service 151 and firewall services 114, 115 and 116 are implemented on a second processor core 161. Tunnel services 150 and 151 may, for example, be implemented as software processes running on the respective processor cores. The firewall

services may be implemented as software threads run by the respective software process (i.e., the respective tunnel services).

An advantage of running separate tunnel services on each processor core is that inter-process communication between the firewall services and communication between a firewall service and the tunnel service is limited to the same core. As a result, when the gateway 100 scales up by adding more processor cores, there will be no loss in performance due to increased inter-process communication.

FIG. 3 shows an exemplary process that may be executed by components of the present disclosure including the gateway 100 according to FIG. 1. By the process, the client connections and, thus, the networking tunnels may be distributed over the different cores 160, 161. The process starts upon establishment of the second connection in step 305 as set out above with reference to step 205 of FIG. 2. The process then proceeds to step 306 where the tunnel module selects one of the tunnel services 150, 151 to which it will forward the second connection. This may be done in several ways. For example, the tunnel module may distribute the connection in a round-robin way by forwarding the connection each time to the next processor core. Alternatively, it may forward the connection to the processor which has the most resources available such as, for example, processing power, memory or any combination thereof. When the processor core is selected, the tunnel module 101 forwards the second connection to the selected tunnel service under step 307, as also set out above with reference to step 206 of FIG. 2.

FIG. 4 illustrates details of an exemplary firewall service 411 that may be implemented in a gateway 100 according to the present disclosure, for example as the firewall service 111 to 116 of FIG. 1. For the sake of exemplary illustration, it will be assumed that firewall service 411 corresponds to firewall service 111, which provides the networking tunnel 131 between client 121 and private network 140. When firewall service 411 is started, it comprises a firewall component 420 for selectively blocking and allowing network traffic between the respective client device 121 and the network devices 141-143 in the private network 140. The rules used by the firewall 420 according to which the traffic is blocked or allowed are stored as firewall rules 421. These rules may, for example, be stored in a file or database residing in memory of the computing system implementing the gateway 100.

Firewall 420 may be a stateful firewall configured to perform stateful packet inspection, thereby keeping track of the state of networking connections established over the networking tunnel 131 between client 121 and network devices 141-143 in private network 140. Each such connection may relate to TCP or UDP networking connections. This way, a finer-grained control is possible as the rules become dependent on the network connection and are thus dynamic. The firewall keeps track of all open network connections by maintaining a list, the state list 422, with all the pending connections. These states may, for example, be stored in a file or database residing in memory of the computing system implementing the gateway 100.

Each firewall service 411 thus maintains a separate state list 422. An advantage of this is that the size of the state list of the firewall service 411 does not grow with the amount of clients connected to the gateway 100. This way, an increase in connected clients does not result in a performance loss due to an increase in the amount of states that have to be managed by a single firewall service.

Firewall service **411** may further include an encryption and/or decryption component **423** for respectively encrypting and decrypting data transmitted to or received from the client **121**. Encryption and decryption may further be performed according to the Internet Protocol Security (or “IPsec protocol,”) Transport Layer Security (or “TLS”) and/or Datagram Transport Layer Security (or “DTLS”). Encrypting and decrypting may further be hardware accelerated by a hardware component in the processor core on which the firewall service **411** is running.

FIG. **5** shows an exemplary process that may be executed by components of the present disclosure, for example by the firewall service **411** of FIG. **4**. By this process, a backup of the firewall service is made such that it can be used during a failure of the computer system, and thus the gateway on which the firewall service **411** is running. In step **501**, firewall **420** allows the establishment of a network connection between client **121** and a networking device or networking application in the private network **140**. This network connection is thus established over the already-established networking tunnel **131**. In order to keep track of the traffic over this new connection and, thus, to allow traffic of this connection to pass the firewall **420**, the firewall registers the connection with the firewall state list **422** under step **501**. Then, in step **502**, the firewall **420** copies the state list **422** and signs the copy in step **503** with a secret signature key. This way it can later be verified that the copy was not altered by an unauthorized party. In step **504**, the copy is then transmitted to the client **121**. This way, the state list and thus the state of the firewall **420** is synchronized with the client **121**. As the state list is signed under step **502**, modifications to the state list can be discovered by checking the signature.

By the process of FIG. **5**, a separate backup can be made of each firewall in the gateway **100**. This has several advantages. First of all, this can be done quickly as only a limited set of states need to be transmitted each time. As a result, when the system is scaled, the time for synchronization of a single firewall service **411** does not increase. Second, this allows for a fail-over arrangement which is scalable as will be demonstrated below.

FIG. **6** illustrates a system according to one embodiment for protecting access to a private network **140**. Access is controlled by gateways **600**, **601** and **602** corresponding to gateway **100** by also implementing tunnel module **101** (not shown), management module **102** (not shown), and one or more tunnel services **150**, **151** (not shown) each implementing one or more firewall services. For sake of clarity, only two firewall services **611A** and **616A** are shown in gateway **600**.

FIG. **7** shows an exemplary process that may be executed by the components of FIG. **6** in the case of a failure of one of the gateways **600-602** (in the following example, a failure of gateway **600** is discussed). In step **701** of the process, gateway **600** establishes a first networking tunnel **631** with client device **621** in order to provide the client **621** access to networking devices **141-143** in the private network. By this, the firewall service **611A** is started thereby implementing a stateful firewall and further including the firewall rules and state list. Step **701** may, for example, be performed in accordance with FIG. **2**.

After establishment of the networking tunnel **631**, firewall service **611A** continuously synchronizes the state list with the client **621** by regularly sending a copy or update **650** of the state list to the client device **621**. For example, the state only needs to be synced if the actual state is changed. Thus, if the client does not open new connections to the protected

servers, no state synchronization will occur. Also, this synchronization is configurable and can be done, for example, continuously, or at least every 1-5 seconds, etc. As long as the networking tunnel is in place **631**, the synchronization is active as illustrated by the loop **703-702**. The synchronization of the state list may further be performed according to the process as illustrated in FIG. **5**.

For purposes of exemplary discussion, consider that at a certain moment in time, the networking tunnel **631** is interrupted. As a result of this interruption, the process proceeds to step **704**. The interruption of the networking tunnel **631** may be caused by several factors such as, for example, by a failure in the networking path between the client **621** and the gateway **600**, or by a failure of the gateway **600** itself. When the client detects the failure of the networking tunnel, it establishes a new networking tunnel **632** with second gateway **601**, which also provides access to private network **140**. This may, for example, be done by the process as illustrated by FIG. **2** by forwarding the tunnel authentication information **653** and client access list **652** to the second gateway **601**.

More specifically, in one embodiment, the information about all of the available gateways (**600**, **601**, **602**) is part of the tunnel lists described above. This information may come from, for example, a central authentication server (controller) such as described in U.S. Pat. No. 9,148,408 B1, which was incorporated by reference above. This information is available at the first connection (i.e., before the first connection may fail). The central authentication server simply picks another candidate from the tunnel list and establishes the connection. The only source information coming from copy or update **650** is the actual connection state from the firewall service that will be synced once the new tunnel is established.

Upon establishment of the tunnel **632**, the firewall service **611B** is started on gateway **601** with identical firewall rules as used in firewall service **611A** because the same access list **652** was used by the management module of gateway **601** to derive the firewall rules.

Upon establishment of the second network tunnel **632**, the client **621** also transmits the state list **651** to the gateway **601** under step **705** of the process. In one embodiment, the copy or update **650** is the source of state list **651**. The connections were going to tunnel **631** before the failover. During the active connections, the state has been continuously synced to client **621**, and also updated when the states are updated for firewall service **611A**. When the tunnel is dropped, a new tunnel is established and the latest received state **651** is synced to firewall service **611B**.

The state list **651** may, for example, be forwarded to the management module. The management module checks that the state list was not changed by the client, and then provides the state list to the firewall service **611B** that uses it as its state list in order to restore the state of the firewall under step **706**. After restoring the state of the firewall, the client resumes, under step **707**, its communication with the private network **140**, thereby achieving a continuous fail-over (i.e., without a loss of the pending network connections established over the initial networking tunnel **631**).

Alternatively, the state list may also be sent over the networking tunnel directly to the firewall service **611B**, wherein this service will check that the state list was not altered. This checking may be done by using the same (shared) secret signature key between the gateways. In one embodiment, because the firewall states are temporarily stored in memory on the client, there is a potential risk of unauthorized tampering with the states. In order to avoid this

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risk, the states are signed with a secret shared key between the gateways. So, if the states are tampered with, the signature is broken and the states will not be accepted.

The process of FIG. 7 has the advantage that no synchronization of firewall rules and states between the gateways **600** and **601** is needed. Therefore, the system can be easily scaled by adding a gateway (e.g., gateway **602**) without bandwidth requirements between the gateways (note that there remains no requirements for any bandwidth to be utilized between the gateways). Furthermore, upon failure of gateway **600**, the networking tunnels can be transferred across the gateways **601** and **602**. For example, the firewall service **616A** may be moved to gateway **602** as service **616B** while firewall **616A** is transferred to gateway **601** as service **611B**.

FIG. 8 shows a suitable exemplary computing system **800** for implementing the gateway according to the above embodiments. Computing system **800** may in general be formed as a suitable general purpose computer and comprise a bus **810**, one or more processor cores **802**, a local memory **804**, one or more optional input interfaces **814**, one or more optional output interfaces **816**, one or more communication interfaces **812**, a storage element interface **806** and one or more storage elements **808**. Bus **810** may comprise one or more conductors that permit communication among the components of the computing system **800**. Processor cores **802** may include any type of conventional processor or microprocessor that interprets and executes programming instructions. Local memory **804** may include a random access memory (RAM) or another type of dynamic storage device that stores information and instructions for execution by processor cores **802** and/or a read only memory (ROM) or another type of static storage device that stores static information and instructions for use by processor **802**. Input interface **814** may comprise one or more conventional mechanisms that permit an operator to input information to the computing device **800**, such as a keyboard **820**, a mouse **830**, a pen, voice recognition and/or biometric mechanisms, etc. Output interface **816** may comprise one or more conventional mechanisms that output information to the operator, such as a display **840**. Communication interface **812** may comprise any transceiver-like mechanism such as, for example, one or more Ethernet interfaces that enables computing system **800** to communicate with other devices and/or systems **801**. The communication interface **812** of computing system **800** may be connected to such another computing system by means of a local area network (LAN) or a wide area network (WAN) such as, for example, the internet. Storage element interface **806** may comprise a storage interface such as, for example, a Serial Advanced Technology Attachment (SATA) interface or a Small Computer System Interface (SCSI) for connecting bus **810** to one or more storage elements **808**, such as one or more local disks, for example SATA disk drives, and control the reading and writing of data to and/or from these storage elements **808**. Although the storage elements **808** above are described as, for example, a local disk, in general any other suitable computer-readable media such as a removable magnetic disk, optical storage media such as a CD-ROM or DVD-ROM disks, solid state drives, flash memory cards, etc., could be used. The system **800** described above may also run as a virtual machine above the physical hardware.

Steps performed according to the above processes may be implemented as computer-executable instructions. These instructions may then be executed on processor cores **802** upon performing the processes. This way, steps executed by tunnel module **101**, management module **102**, tunnel ser-

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vices **150**, **151** and firewall services **111-116** may, for example, be implemented as instructions on computing system **800** thereby implementing gateway **100**. Processor cores **802** may correspond to processor cores **160**, **161** of gateway **100**. A tunnel service **150**, **151** is then run on each of the processor cores **802**. Each processor core **802** then runs separate firewall services respectively started by one of the tunnel services. Data packet communication between client devices **121-126** and the gateway **100** may be performed over a networking interface **812**. Also, data packets communicated between gateway **100** and private network **140** may be exchanged over a networking interface **812**. The computer-executable instructions may form or be part of a computer program product that is stored on storage element **808** or any computer readable storage medium.

Communication among systems, devices, and components operating in conjunction with embodiments of the present disclosure may be performed using any suitable communication method, such as, for example, a telephone network, an extranet, an intranet, the Internet, point of interaction device (point of sale device, personal digital assistant (e.g., iPhone®, Palm Pilot®, BlackBerry®), cellular phone, kiosk, etc.), online communications, satellite communications, off-line communications, wireless communications, transponder communications, local area network (LAN), wide area network (WAN), virtual private network (VPN), networked or linked devices, keyboard, mouse and/or any suitable communication or data input modality. Systems and devices of the present disclosure may utilize TCP/IP communications protocols as well as IPX, Appletalk, IP-6, NetBIOS, OSI, any tunnelling protocol (e.g. IPsec, SSH), or any number of existing or future protocols.

While some embodiments can be implemented in fully functioning computers and computer systems, various embodiments are capable of being distributed as a computing product in a variety of forms and are capable of being applied regardless of the particular type of machine or computer-readable media used to actually effect the distribution.

A machine readable medium can be used to store software and data which when executed by a data processing system causes the system to perform various methods. The executable software and data may be stored in various places including, for example, ROM, volatile RAM, non-volatile memory and/or cache. Portions of this software and/or data may be stored in any one of these storage devices. Further, the data and instructions can be obtained from centralized servers or peer to peer networks. Different portions of the data and instructions can be obtained from different centralized servers and/or peer to peer networks at different times and in different communication sessions or in a same communication session. The data and instructions can be obtained in entirety prior to the execution of the applications. Alternatively, portions of the data and instructions can be obtained dynamically, just in time, when needed for execution. Thus, it is not required that the data and instructions be on a machine readable medium in entirety at a particular instance of time.

Examples of computer-readable media include but are not limited to recordable and non-recordable type media such as volatile and non-volatile memory devices, read only memory (ROM), random access memory (RAM), flash memory devices, floppy and other removable disks, magnetic disk storage media, optical storage media (e.g., Compact Disk Read-Only Memory (CD ROMS), Digital Versatile Disks (DVDs), etc.), among others. The computer-readable media may store the instructions.

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In various embodiments, hardwired circuitry may be used in combination with software instructions to implement the techniques. Thus, the techniques are neither limited to any specific combination of hardware circuitry and software nor to any particular source for the instructions executed by the data processing system.

Although some of the drawings illustrate a number of operations in a particular order, operations which are not order dependent may be reordered and other operations may be combined or broken out. While some reordering or other groupings are specifically mentioned, others will be apparent to those of ordinary skill in the art and so do not present an exhaustive list of alternatives. Moreover, it should be recognized that the stages could be implemented in hardware, firmware, software or any combination thereof.

For the sake of brevity, conventional data networking, application development and other functional aspects of the systems (and components of the individual operating components of the systems) may not be described in detail herein. Furthermore, the connecting lines shown in the various figures contained herein are intended to represent exemplary functional relationships and/or physical couplings between the various elements. It should be noted that many alternative or additional functional relationships or physical connections may be present in a practical system.

The various system components discussed herein may include one or more of the following: a host server or other computing systems including a processor for processing digital data; a memory coupled to the processor for storing digital data; an input digitizer coupled to the processor for inputting digital data; an application program stored in the memory and accessible by the processor for directing processing of digital data by the processor; a display device coupled to the processor and memory for displaying information derived from digital data processed by the processor; and a plurality of databases. Various databases used herein may include: shipping data, package data, and/or any data useful in the operation of the system.

Various functionality may be performed via a web browser and/or application interfacing utilizing a web browser. Such browser applications may comprise Internet browsing software installed within a computing unit or a system to perform various functions. These computing units or systems may take the form of a computer or set of computers, and any type of computing device or systems may be used, including laptops, notebooks, tablets, handheld computers, personal digital assistants, set-top boxes, workstations, computer-servers, main frame computers, mini-computers, PC servers, network sets of computers, personal computers and tablet computers, such as iPads, iMACs, and MacBook, kiosks, terminals, point of sale (POS) devices and/or terminals, televisions, or any other device capable of receiving data over a network. Various embodiments may utilize Microsoft Internet Explorer, Mozilla Firefox, Google Chrome, Apple Safari, Opera, or any other of the myriad software packages available for browsing the internet.

Various embodiments may operate in conjunction with any suitable operating system (e.g., Windows NT, 95/98/2000/CE/Mobile/, Windows 7/8, OS2, UNIX, Linux, Solaris, MacOS, PalmOS, etc.) as well as various conventional support software and drivers typically associated with computers. Various embodiments may include any suitable personal computer, network computer, workstation, personal digital assistant, cellular phone, smart phone, minicomputer, mainframe or the like. Embodiments may implement security protocols, such as Secure Sockets Layer (SSL), Trans-

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port Layer Security (TLS), and Secure Shell (SSH). Embodiments may implement any desired application layer protocol, including http, https, ftp, and sftp.

The various system components may be independently, separately or collectively suitably coupled to a network via data links which includes, for example, a connection to an Internet Service Provider (ISP) over the local loop as is typically used in connection with standard modem communication, cable modem, satellite networks, ISDN, Digital Subscriber Line (DSL), or various wireless communication methods. It is noted that embodiments of the present disclosure may operate in conjunction with any suitable type of network, such as an interactive television (ITV) network.

The system may be partially or fully implemented using cloud computing. "Cloud" or "Cloud computing" includes a model for enabling convenient, on-demand network access to a shared pool of configurable computing resources (e.g., networks, servers, storage, applications, and services) that can be rapidly provisioned and released with minimal management effort or service provider interaction. Cloud computing may include location-independent computing, whereby shared servers provide resources, software, and data to computers and other devices on demand.

Various embodiments may be used in conjunction with web services, utility computing, pervasive and individualized computing, security and identity solutions, autonomic computing, cloud computing, commodity computing, mobility and wireless solutions, open source, biometrics, grid computing and/or mesh computing.

Any databases discussed herein may include relational, hierarchical, graphical, or object-oriented structure and/or any other database configurations. Moreover, the databases may be organized in any suitable manner, for example, as data tables or lookup tables. Each record may be a single file, a series of files, a linked series of data fields or any other data structure. Association of certain data may be accomplished through any desired data association technique such as those known or practiced in the art. For example, the association may be accomplished either manually or automatically.

Any databases, systems, devices, servers or other components of the system may be located at a single location or at multiple locations, wherein each database or system includes any of various suitable security features, such as firewalls, access codes, encryption, decryption, compression, decompression, and/or the like.

Encryption may be performed by way of any of the techniques now available in the art or which may become available—e.g., Twofish, RSA, El Gamal, Schorr signature, DSA, PGP, PKI, and symmetric and asymmetric cryptosystems.

Embodiments may connect to the Internet or an intranet using standard dial-up, cable, DSL or any other Internet protocol known in the art. Transactions may pass through a firewall in order to prevent unauthorized access from users of other networks.

The computers discussed herein may provide a suitable website or other Internet-based graphical user interface which is accessible by users. For example, the Microsoft Internet Information Server (IIS), Microsoft Transaction Server (MTS), and Microsoft SQL Server, may be used in conjunction with the Microsoft operating system, Microsoft NT web server software, a Microsoft SQL Server database system, and a Microsoft Commerce Server. Additionally, components such as Access or Microsoft SQL Server, Oracle, Sybase, Informix MySQL, Interbase, etc., may be used to provide an Active Data Object (ADO) compliant database management system. In another example, an

Apache web server can be used in conjunction with a Linux operating system, a MySQL database, and the Perl, PHP, and/or Python programming languages.

Any of the communications, inputs, storage, databases or displays discussed herein may be facilitated through a website having web pages. The term "web page" as it is used herein is not meant to limit the type of documents and applications that might be used to interact with the user. For example, a typical website might include, in addition to standard HTML documents, various forms, Java applets, JavaScript, active server pages (ASP), common gateway interface scripts (CGI), extensible markup language (XML), dynamic HTML, cascading style sheets (CSS), AJAX (Asynchronous Javascript And XML), helper applications, plug-ins, and the like. A server may include a web service that receives a request from a web server, the request including a URL and an IP address. The web server retrieves the appropriate web pages and sends the data or applications for the web pages to the IP address. Web services are applications that are capable of interacting with other applications over a communications means, such as the Internet.

Various embodiments may employ any desired number of methods for displaying data within a browser-based document. For example, data may be represented as standard text or within a fixed list, scrollable list, drop-down list, editable text field, fixed text field, pop-up window, and the like. Likewise, embodiments may utilize any desired number of methods for modifying data in a web page such as, for example, free text entry using a keyboard, selection of menu items, check boxes, option boxes, and the like.

The exemplary systems and methods illustrated herein may be described in terms of functional block components, screen shots, optional selections and various processing steps. It should be appreciated that such functional blocks may be realized by any number of hardware and/or software components configured to perform the specified functions. For example, the system may employ various integrated circuit components, e.g., memory elements, processing elements, logic elements, look-up tables, and the like, which may carry out a variety of functions under the control of one or more microprocessors or other control devices. Similarly, the software elements of the system may be implemented with any programming or scripting language such as C, C++, C#, Java, JavaScript, VBScript, Macromedia Cold Fusion, COBOL, Microsoft Active Server Pages, assembly, PERL, PHP, AWK, Python, Visual Basic, SQL Stored Procedures, PL/SQL, any UNIX shell script, and extensible markup language (XML) with the various algorithms being implemented with any combination of data structures, objects, processes, routines or other programming elements. Further, it should be noted that the system may employ any number of conventional techniques for data transmission, signalling, data processing, network control, and the like. Still further, the system could be used to detect or prevent security issues with a client-side scripting language, such as JavaScript, VBScript or the like.

The systems and methods of the present disclosure may be embodied as a customization of an existing system, an add-on product, a processing apparatus executing upgraded software, a stand alone system, a distributed system, a method, a data processing system, a device for data processing, and/or a computer program product. Accordingly, any portion of the system or a module may take the form of a processing apparatus executing code, an internet based embodiment, an entirely hardware embodiment, or an embodiment combining aspects of the internet, software and hardware. Furthermore, the system may take the form of a

computer program product on a computer-readable storage medium having computer-readable program code means embodied in the storage medium. Any suitable computer-readable storage medium may be utilized, including hard disks, CD-ROM, optical storage devices, magnetic storage devices, and/or the like.

The system and method is described herein with reference to screen shots, block diagrams and flowchart illustrations of methods, apparatus (e.g., systems), and computer program products according to various embodiments. It will be understood that each functional block of the block diagrams and the flowchart illustrations, and combinations of functional blocks in the block diagrams and flowchart illustrations, respectively, can be implemented by computer program instructions.

These computer program instructions may be loaded onto a general purpose computer, special purpose computer, or other programmable data processing apparatus to produce a machine, such that the instructions that execute on the computer or other programmable data processing apparatus create means for implementing the functions specified in the flowchart block or blocks. These computer program instructions may also be stored in a computer-readable memory that can direct a computer or other programmable data processing apparatus to function in a particular manner, such that the instructions stored in the computer-readable memory produce an article of manufacture including instruction means which implement the function specified in the flowchart block or blocks. The computer program instructions may also be loaded onto a computer or other programmable data processing apparatus to cause a series of operational steps to be performed on the computer or other programmable apparatus to produce a computer-implemented process such that the instructions which execute on the computer or other programmable apparatus provide steps for implementing the functions specified in the flowchart block or blocks.

Accordingly, functional blocks of the block diagrams and flowchart illustrations support combinations of means for performing the specified functions, combinations of steps for performing the specified functions, and program instruction means for performing the specified functions. It will also be understood that each functional block of the block diagrams and flowchart illustrations, and combinations of functional blocks in the block diagrams and flowchart illustrations, can be implemented by either special purpose hardware-based computer systems which perform the specified functions or steps, or suitable combinations of special purpose hardware and computer instructions. Further, illustrations of the process flows and the descriptions thereof may make reference to user windows, webpages, websites, web forms, prompts, etc. Practitioners will appreciate that the illustrated steps described herein may comprise in any number of configurations including the use of windows, webpages, web forms, popup windows, prompts and the like. It should be further appreciated that the multiple steps as illustrated and described may be combined into single webpages and/or windows but have been expanded for the sake of simplicity. In other cases, steps illustrated and described as single process steps may be separated into multiple webpages and/or windows but have been combined for simplicity.

The term "non-transitory" is to be understood to remove only propagating transitory signals per se from the claim scope and does not relinquish rights to all standard computer-readable media that are not only propagating transitory signals per se. Stated another way, the meaning of the term

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“non-transitory computer-readable medium” should be construed to exclude only those types of transitory computer-readable media which were found in *In Re Nuijten* to fall outside the scope of patentable subject matter under 35 U.S.C. §101.

Benefits, other advantages, and solutions to problems have been described herein with regard to specific embodiments. However, the benefits, advantages, solutions to problems, and any elements that may cause any benefit, advantage, or solution to occur or become more pronounced are not to be construed as critical, required, or essential features or elements of the disclosure.

Although the disclosure includes a method, it is contemplated that it may be embodied as computer program instructions on a tangible computer-readable carrier, such as a magnetic or optical memory or a magnetic or optical disk. All structural, chemical, and functional equivalents to the elements of the above-described exemplary embodiments that are known to those of ordinary skill in the art are expressly incorporated herein by reference and are intended to be encompassed by the present claims. Moreover, it is not necessary for a device or method to address each and every problem sought to be solved by the present disclosure, for it to be encompassed by the present claims. Furthermore, no element, component, or method step in the present disclosure is intended to be dedicated to the public regardless of whether the element, component, or method step is explicitly recited in the claims. No claim element herein is to be construed under the provisions of 35 U.S.C. 112, sixth paragraph, unless the element is expressly recited using the phrase “means for.” As used herein, the terms “comprises”, “comprising”, or any other variation thereof, are intended to cover a non-exclusive inclusion, such that a process, method, article, or apparatus that comprises a list of elements does not include only those elements but may include other elements not expressly listed or inherent to such process, method, article, or apparatus.

Where a phrase similar to “at least one of A, B, or C,” “at least one of A, B, and C,” “one or more A, B, or C,” or “one or more of A, B, and C” is used, it is intended that the phrase be interpreted to mean that A alone may be present in an embodiment, B alone may be present in an embodiment, C alone may be present in an embodiment, or that any combination of the elements A, B and C may be present in a single embodiment; for example, A and B, A and C, B and C, or A and B and C.

Changes and modifications may be made to the disclosed embodiments without departing from the scope of the present disclosure. These and other changes or modifications are intended to be included within the scope of the present disclosure, as expressed in the following claims.

The invention claimed is:

1. A computer-implemented method, comprising:

in response to a request from a client device, establishing, by a first computer system implementing a first gateway to a private network, a first network tunnel between the client device and the first gateway, wherein the private network comprises one or more network devices;

receiving, by the first computer system, from the client device, a client access list indicating those of the network devices in the private network that are allowed to communicate with the client device; and

starting, for the first network tunnel, a separate firewall service with a separate set of firewall rules on the first computer system for selectively blocking and allowing network traffic between the client device and the one or

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more network devices in the private network, wherein each of the firewall rules is derived from the client access list.

2. The method of claim 1, wherein the separate firewall service is a stateful firewall comprising a state listing for keeping track of states of network connections between the client device and the network devices in the private network.

3. The method of claim 2, further comprising synchronizing by the first computer system, the state listing with the client device.

4. The method of claim 3, further comprising making, by the first computer system, the state listing unalterable by the client device.

5. The method of claim 3, further comprising, in response to a failure of the first network tunnel, implementing, by a second computer system, a second gateway to the private network, the implementing comprising:

receiving from the client device a second request to establish a second network tunnel between the client device and the second gateway;

receiving from the client device the state listing;

establishing the second network tunnel between the client device and the second gateway;

after establishing the second network tunnel, starting a second separate firewall service with the separate set of firewall rules on the second computer system for selectively blocking and allowing network traffic between the client device and the one or more network devices in the private network; and

restoring the state of the separate firewall service within the second separate firewall service using the state listing.

6. The method of claim 5, further comprising:

in response to the failure of the first network tunnel, receiving, by the second computer system, from the client device, the client access list; and

determining, by the second computer system, from the client access list, firewall rules for the second separate firewall service.

7. The method of claim 1, wherein:

the first computer system comprises a plurality of processor cores;

the first computer system is configured to run a process on each processor core; and

the starting the separate firewall service further comprises starting, by the process, a new processing thread for providing the separate firewall service on one of the plurality of processor cores.

8. The method of claim 7, wherein the starting the separate firewall service further comprises selecting one of the plurality of processor cores and thereby the process for running the separate firewall service, and the method further comprising, after starting the separate firewall service, passing payload data from the first network tunnel to the new processing thread.

9. The method of claim 7, further comprising:

encrypting, by the new processing thread, payload data for the client device; and

decrypting payload data from the client device.

10. The method of claim 9, wherein the plurality of processor cores comprises hardware acceleration for performing at least one of the encrypting or decrypting.

11. A system, comprising:

at least one processor core, each core configured to run a process; and

memory storing instructions configured to instruct the at least one processor core to:

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in response to a request from a client device, establish a network tunnel between the client device and a gateway to a private network, wherein the private network comprises one or more network devices; receive from the client device, a client access list 5 indicating those of the network devices in the private network that are allowed to communicate with the client device; and

start, for the network tunnel, a separate firewall service with a separate set of firewall rules for selectively blocking and allowing network traffic between the client device and the one or more network devices in the private network, wherein each of the firewall rules is derived from the client access list. 10

12. The system of claim 11, wherein the starting the separate firewall service further comprises starting, by a first process of a first core of the plurality of processor cores, a new processing thread for providing the separate firewall service on a second core of the plurality of processor cores. 15

13. The system of claim 12, wherein the second core runs the separate firewall service. 20

14. The system of claim 13, wherein the instructions are further configured to instruct the at least one processor core to, after starting the separate firewall service, pass payload data from the network tunnel to the new processing thread.

15. The system of claim 11, further comprising memory storing the separate set of firewall rules. 25

16. A non-transitory computer readable storage medium storing computer-readable instructions, which when executed, cause a first computer system to:

in response to a request from a client device, establish a first network tunnel between the client device and a first gateway; 30

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receive, by the first computer system, from the client device, a client access list; and

start, for the first network tunnel, a separate firewall service with a separate set of firewall rules on the first computer system for selectively blocking and allowing network traffic between the client device and one or more network devices in a private network, wherein each of the firewall rules is derived from the client access list.

17. The storage medium of claim 16, wherein the instructions further cause the first computer system to, in response to a failure of the first network tunnel, cause the implementation by a second computer system of a second gateway to the private network, the implementation to establish a second network tunnel between the client device and the second gateway, the implementation further to start a second separate firewall service with the separate set of firewall rules.

18. The storage medium of claim 17, wherein the instructions further cause the first computer system to:

in response to the failure of the first network tunnel, cause providing of the client access list to the second computer system.

19. The storage medium of claim 16, wherein the separate firewall service is a stateful firewall comprising a state listing for keeping track of states of network connections between the client device and the network devices.

20. The storage medium of claim 16, wherein the instructions further cause the first computer system to encrypt payload data for the client device.

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Patente PCT
PROTECCIÓN DE LOS DISPOSITIVOS DE RED POR UN CORTAFUEGOS

Detalle del pago

Id.	Rentístico	Concepto y Servicio	Cant.	Val. Unit.	Subtotal
1	50005-01-01	Solicitud examen de patentabilidad Patente de Invención (DESCUENTO INTERNET)	1	\$ 1,097,600	\$ 1,097,600
TOTAL					\$ 1,097,600